Nearest neighbors: performance

Lecture 7 by Marina Barsky

K-NN algorithm. Summary

- The training set *is the* model
- Advantages:
 - Building a classifier: zero work
 - Updating the model with every new record: zero work
 - Interpretable: we can justify our classification
 - Good for predicting numeric values (Regressor)
- Disadvantages:

– The query is computationally expensive!

K-NN query performance: time and space

Space: O(M x N)

Running time: O(M x N)

- M number of attributes
- N total instances in the training set

K-NN performance improvements

Heuristic algorithms:

- 1. **IB2:** save memory, speed up classification
- 2. IB3: deal with noise

Data structures:

- 1. KD-tree
- 2. Ball-tree

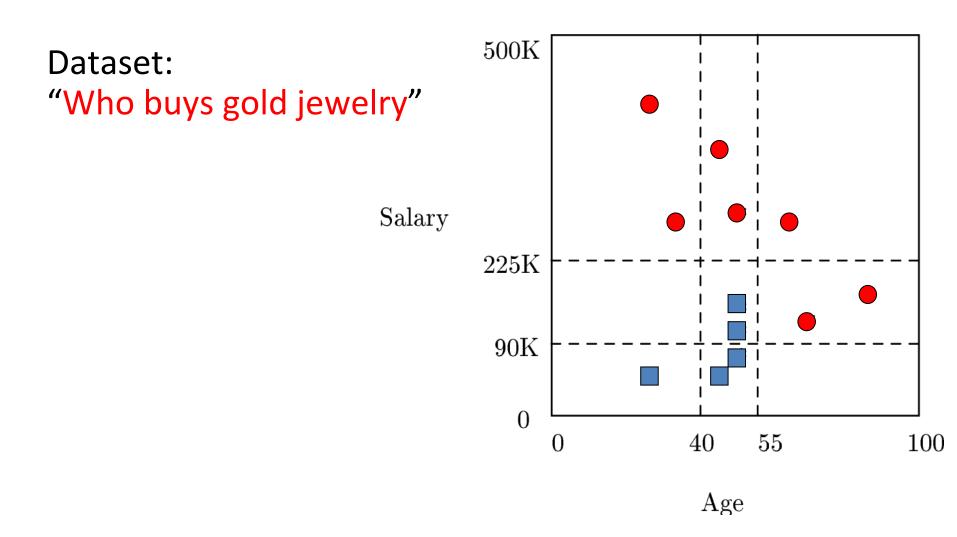
*) IB – Instance Based

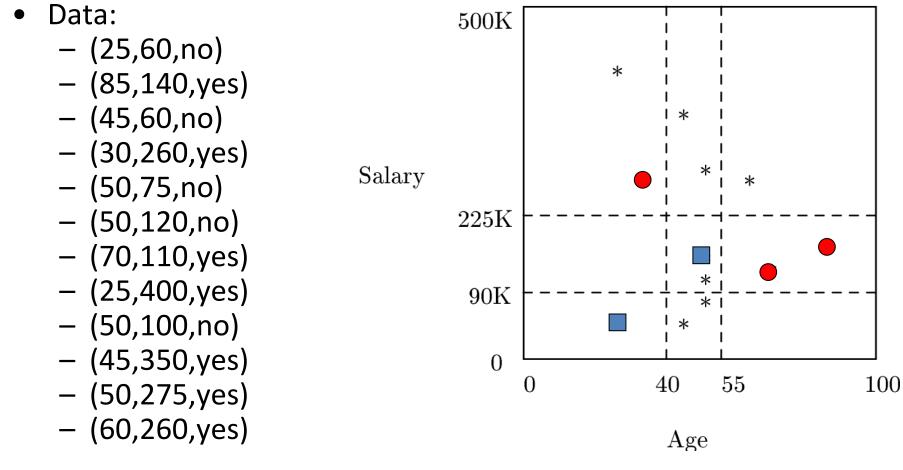
Algorithm example: IB2

main idea

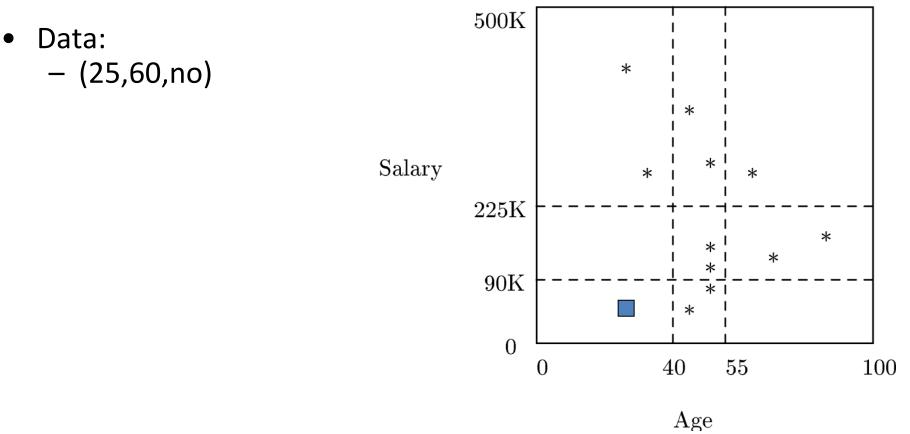
- Work incrementally
- Only incorporate misclassified instances

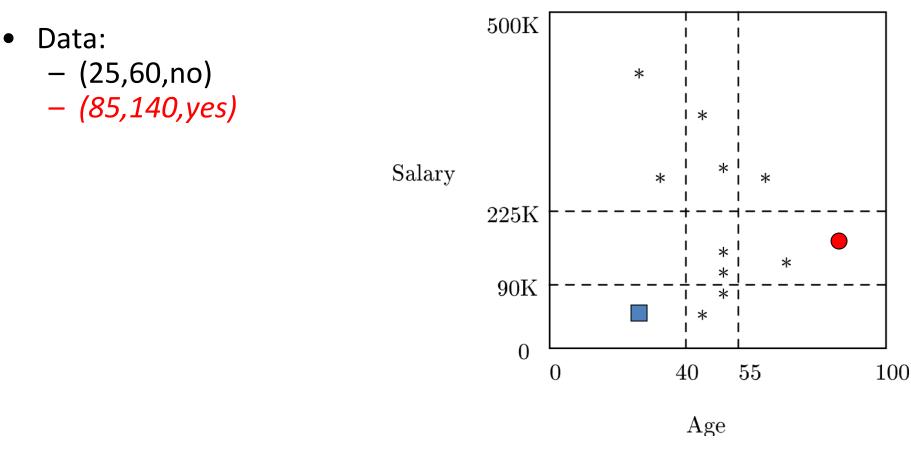
Example: IB2

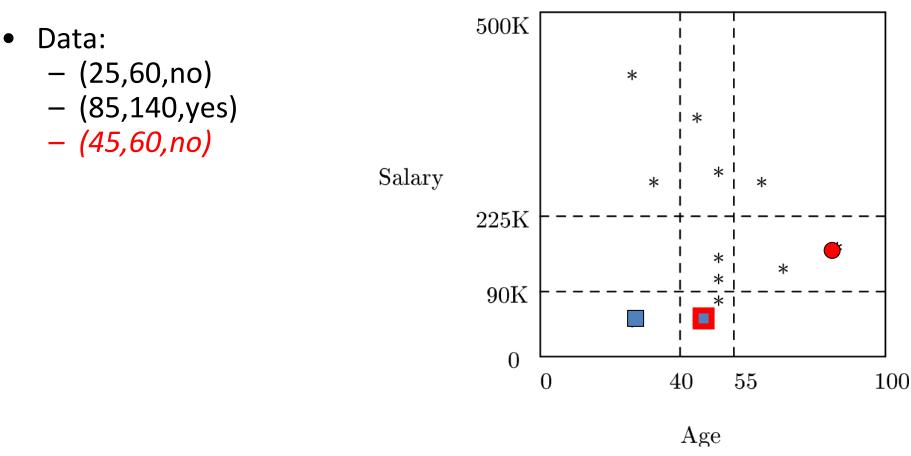


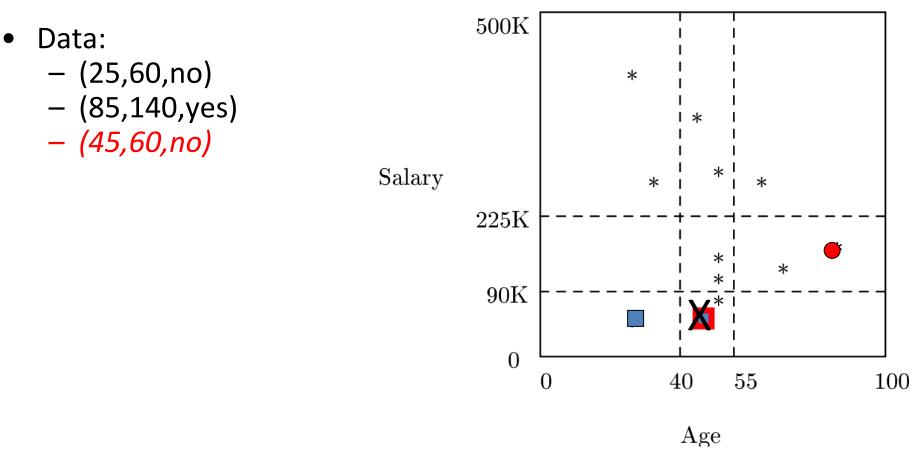


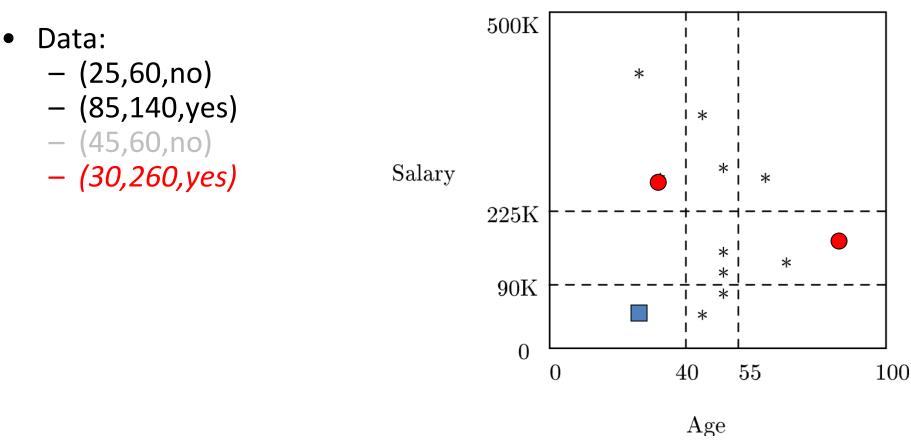
IB2 output: We memorize only these 5 points.

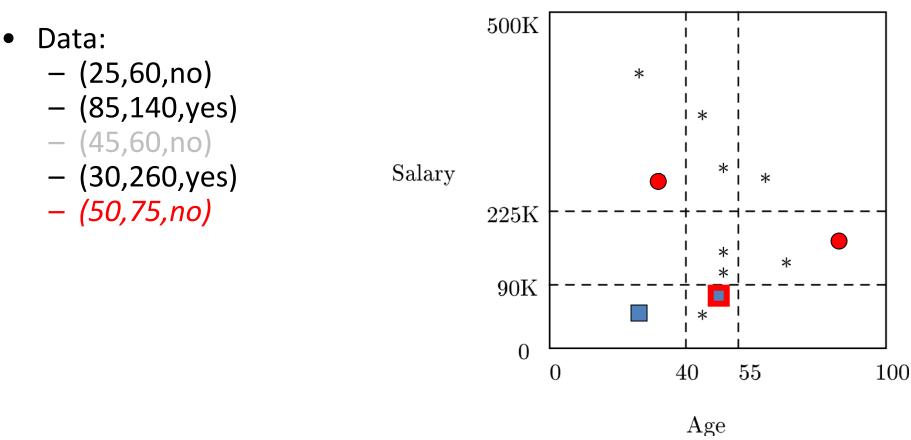


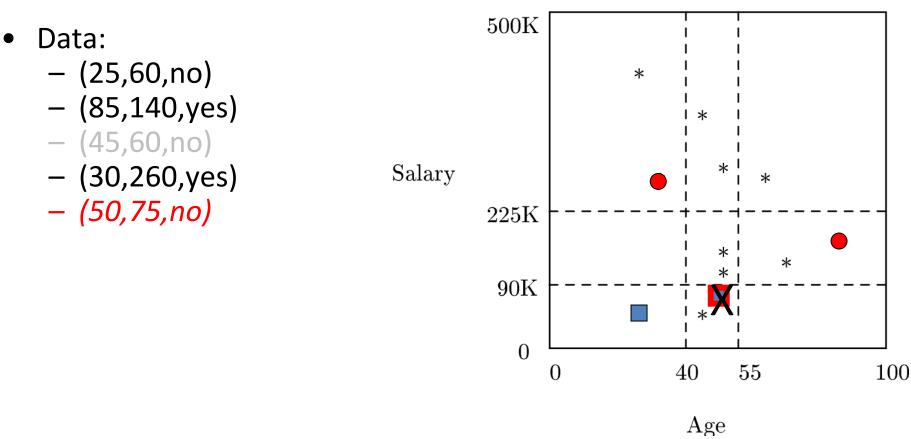


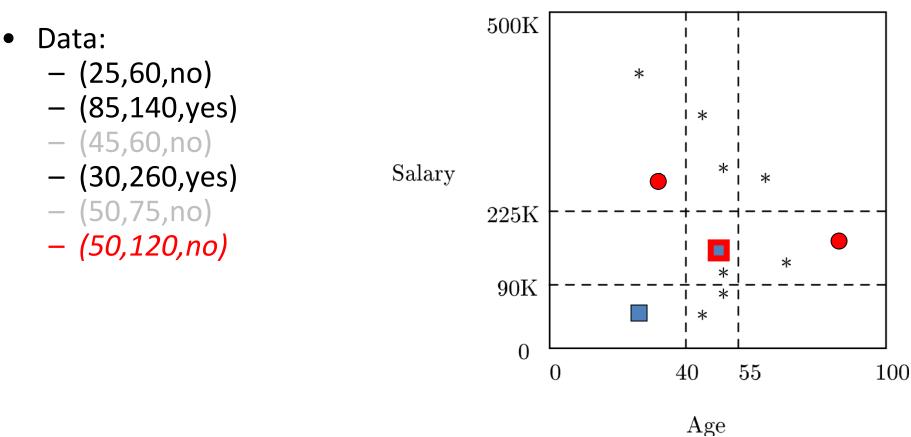






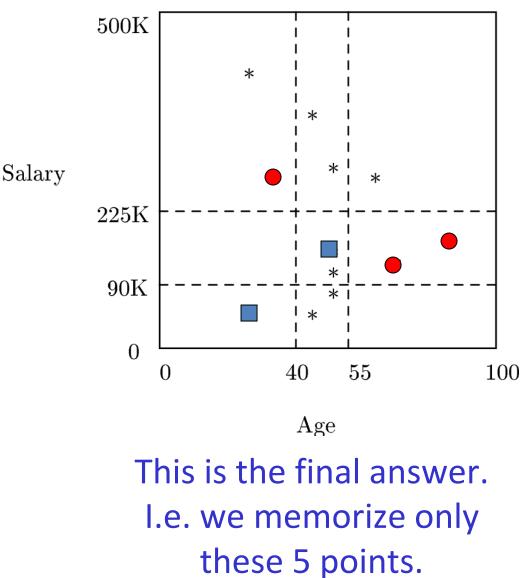








- Continuing in a similar way, we finally get a smaller set to memorize.
 - The colored points are the ones that get memorized.



IB2 summary

- Work incrementally
- Only incorporate misclassified instances
- Problem: noisy data might get incorporated

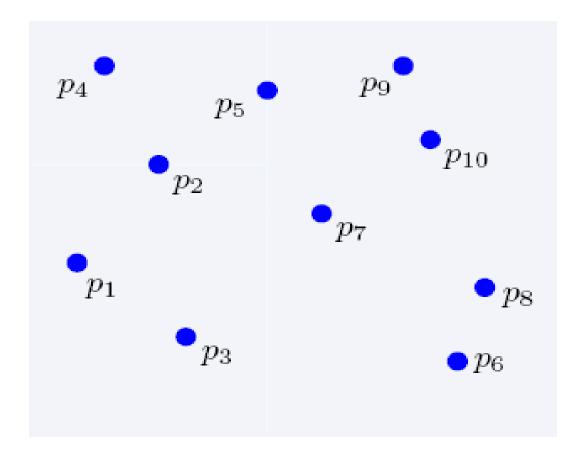
Data structure example: KD-tree

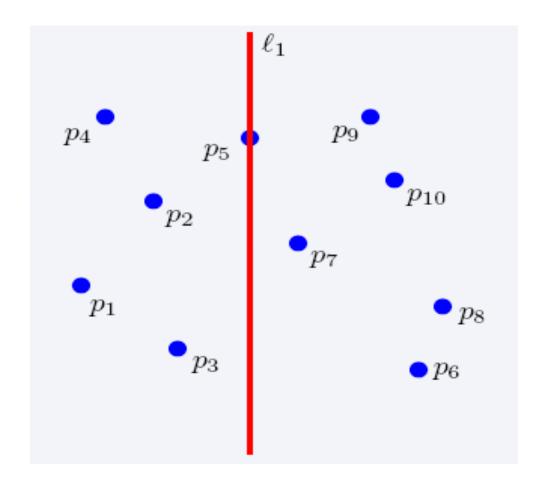
To find nearest neighbors quickly, use a special type of a binary tree: KD-tree

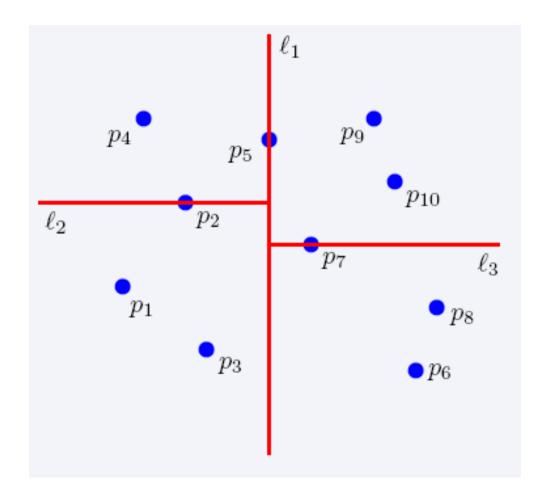
- At the root of the tree we split the set of points into two subsets of same size by a hyperplane vertical to x1-axis (first dimension)
- At the children of the root, the partition is based on the second dimension: x2
- At depth d, we start all over again by partitioning on the first coordinate
- The recursion stops until there is only one point left, which is stored as a leaf

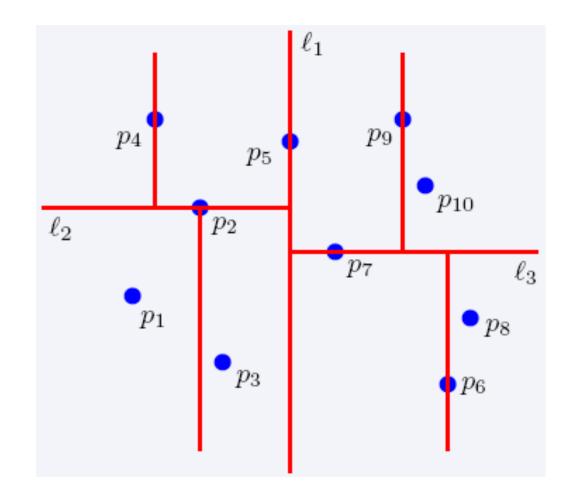
2-dimensional kd-trees

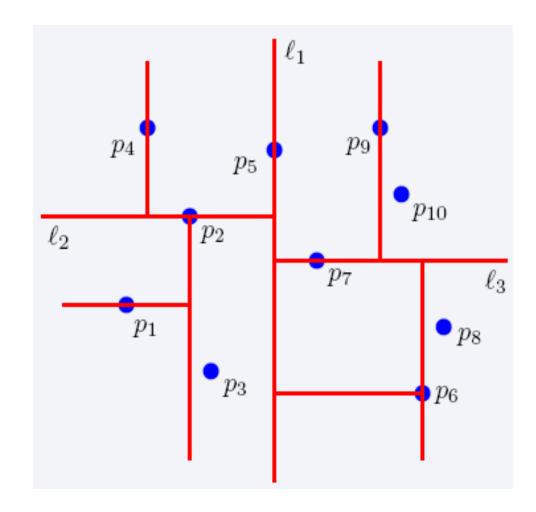
- Algorithm:
 - Choose x or y coordinate (alternate)
 - Choose the median of the coordinate: this defines a horizontal or vertical line
 - Recurse on both sides
- We get a binary tree:
 - Size O(n)
 - Depth O(logn)
 - Construction time O(nlogn)



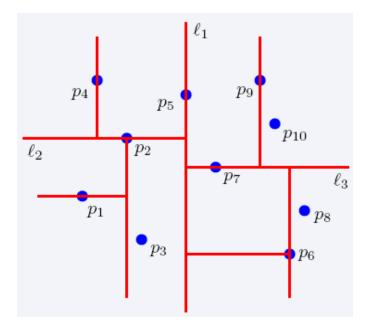


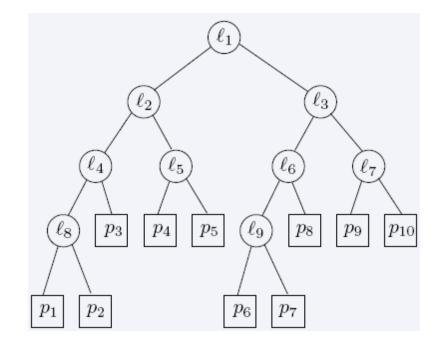




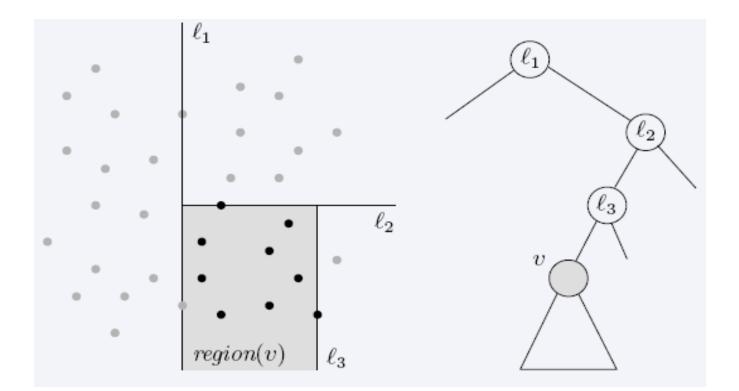


The complete kd-tree





Region of node v



Region(v) : the subtree rooted at **v** stores the points in black dots

d-dimensional kd-trees

- A data structure to support range queries in R^d
- Preprocessing time: O(nlogn)
- Space complexity: O(n)
- Query time: O(n^{1-1/d}+k)